

# Introduction to DNS and its vulnerabilities

Olaf M. Kolkman  
[olaf@nlnetlabs.nl](mailto:olaf@nlnetlabs.nl)

# DNS has a distributed nature

- Authoritative servers all provide part of the name space
- User devices query a local server that maintains a cache
  - For better performance
  - For scalability of the system as a whole

# Terminology

- Authoritative Nameserver: Maintains an authoritative copy of the data.
- Recursive Nameserver: Contacts Authoritative servers to compose an answer for stub resolvers. Also called Caching Nameserver or Cache
- Stub Resolver: fires off queries to pre-configured addresses and expects an answer. Usually implemented in OS library
  - `gethostbyname()`

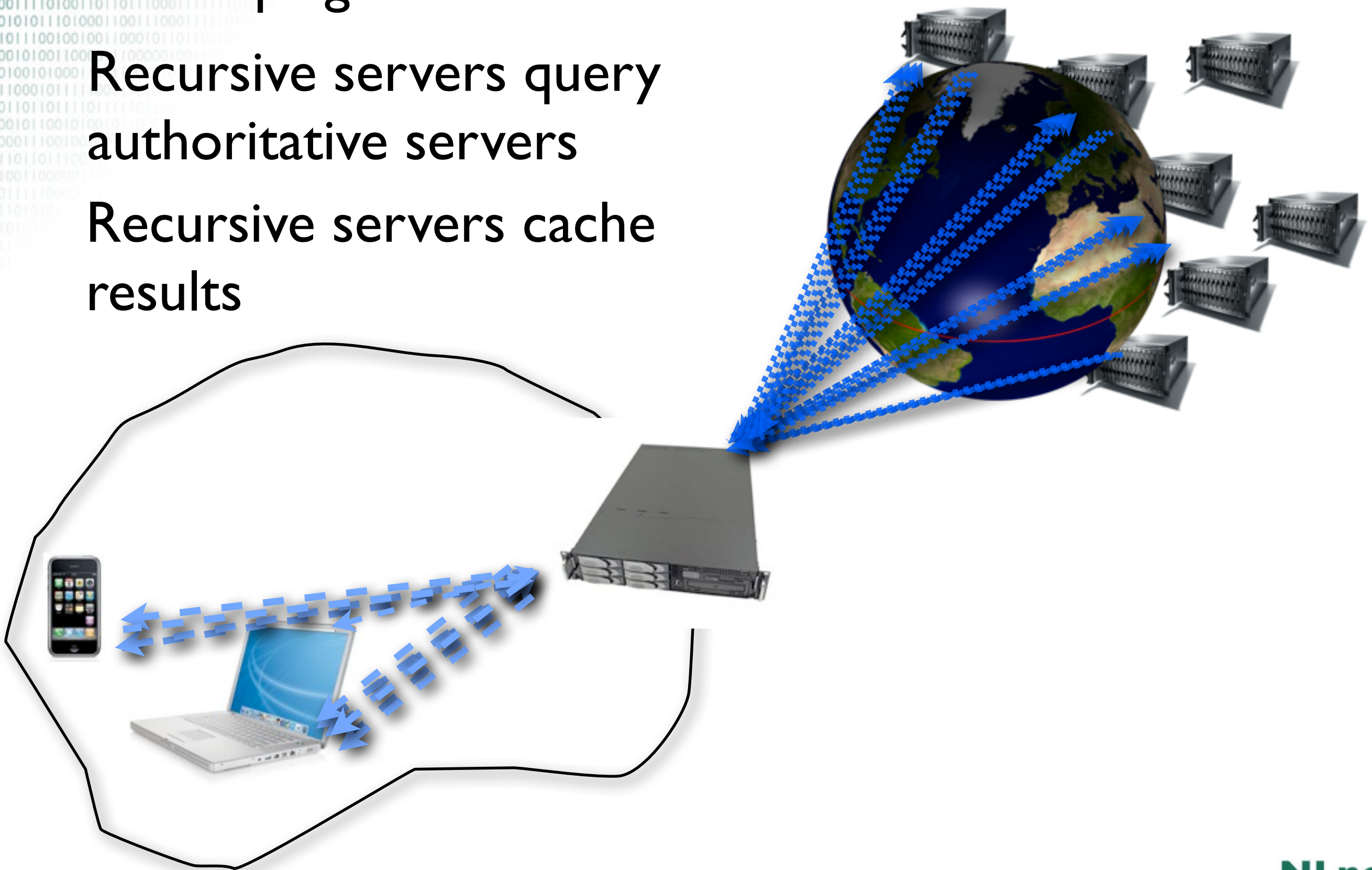
# Animation



Look up against recursive servers

Recursive servers query  
authoritative servers

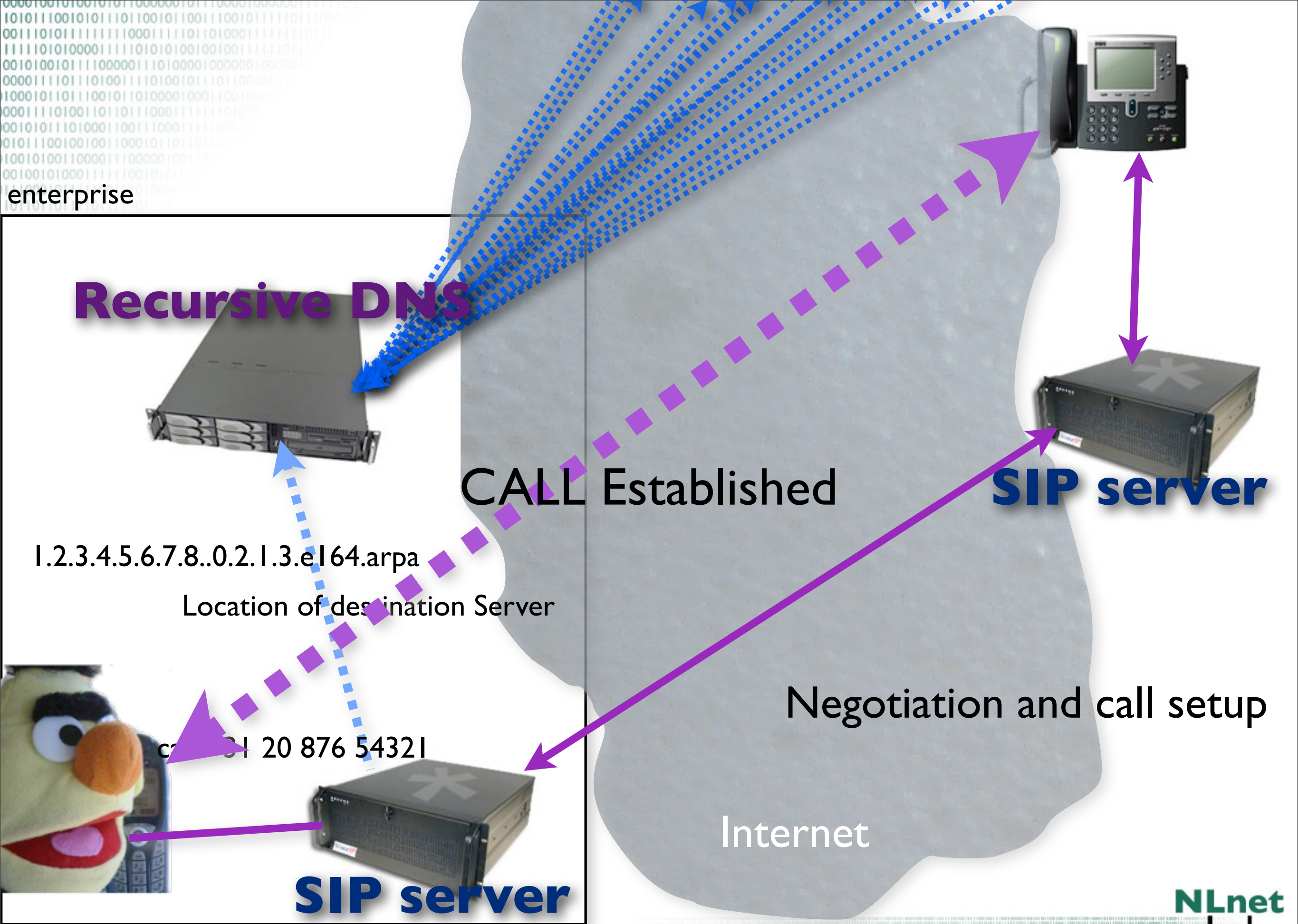
Recursive servers cache  
results



# When do you use the DNS

- Anytime that you need to know where the other guy is
- DNS is the phone book of the Internet
- So it is used when people make a voice over IP call





enterprise

**Recursive DNS**

**CALL Established**

**SIP server**

`1.2.3.4.5.6.7.8..0.2.1.3.e164.arpa`

Location of destination Server

**Negotiation and call setup**

**Internet**

**SIP server**

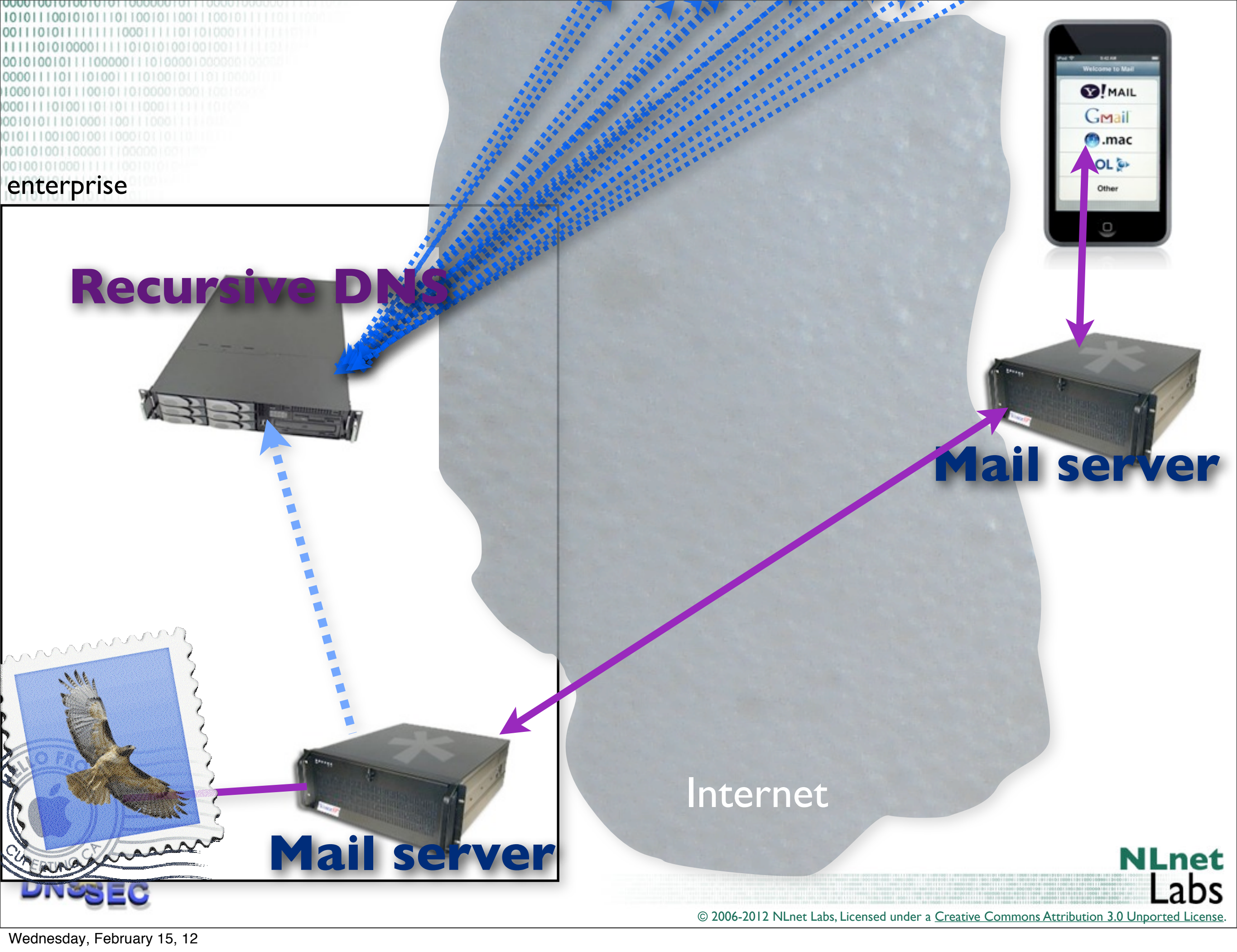
**DNSSEC**

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Or they use the DNS  
when sending MAIL





enterprise

**Recursive DNS**

**Mail server**

Internet

**Mail server**



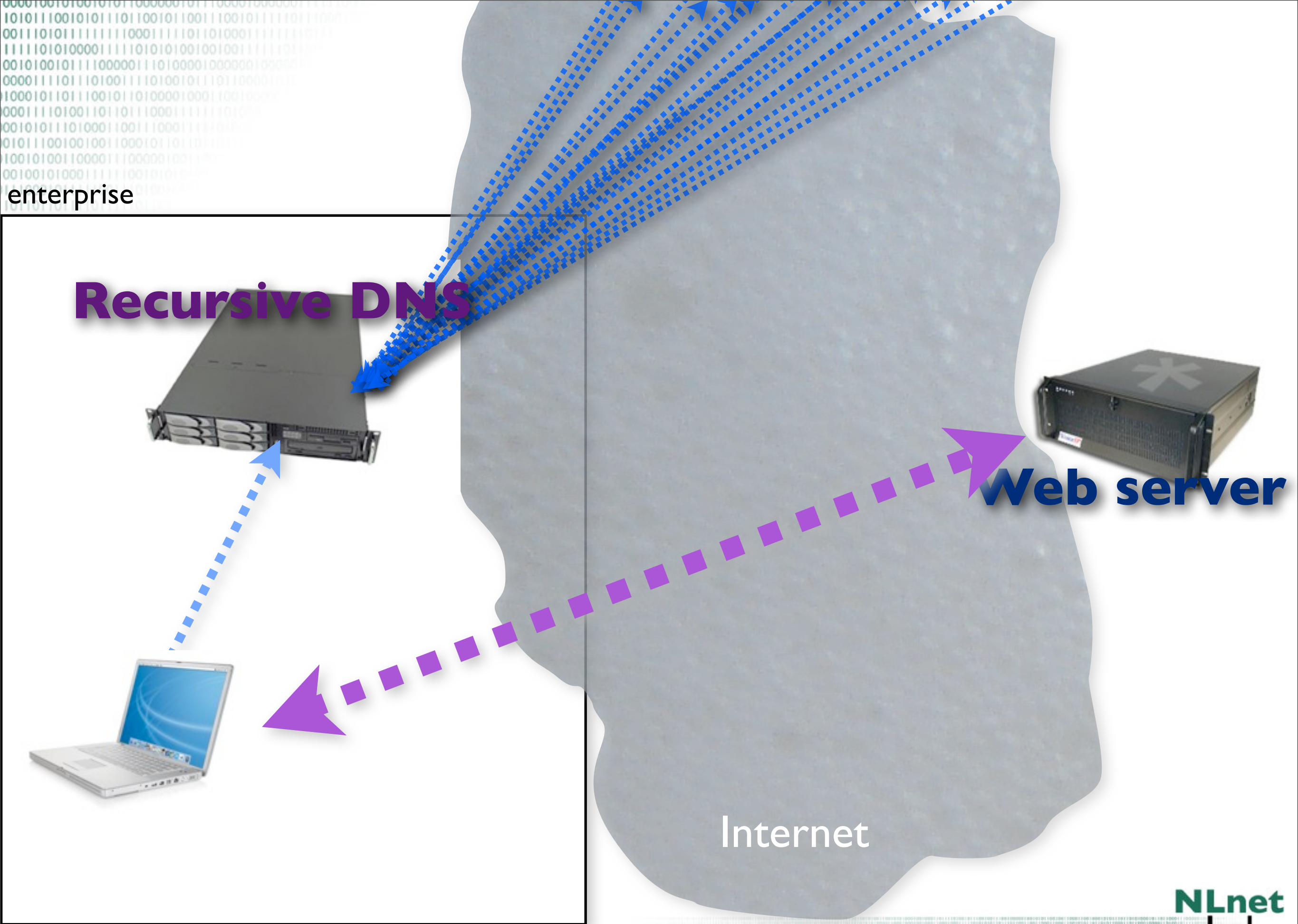
**DNSSEC**

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Or they use the DNS  
when browsing the  
Web





DNSSEC

NLnet  
Labs

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# Or they use the DNS

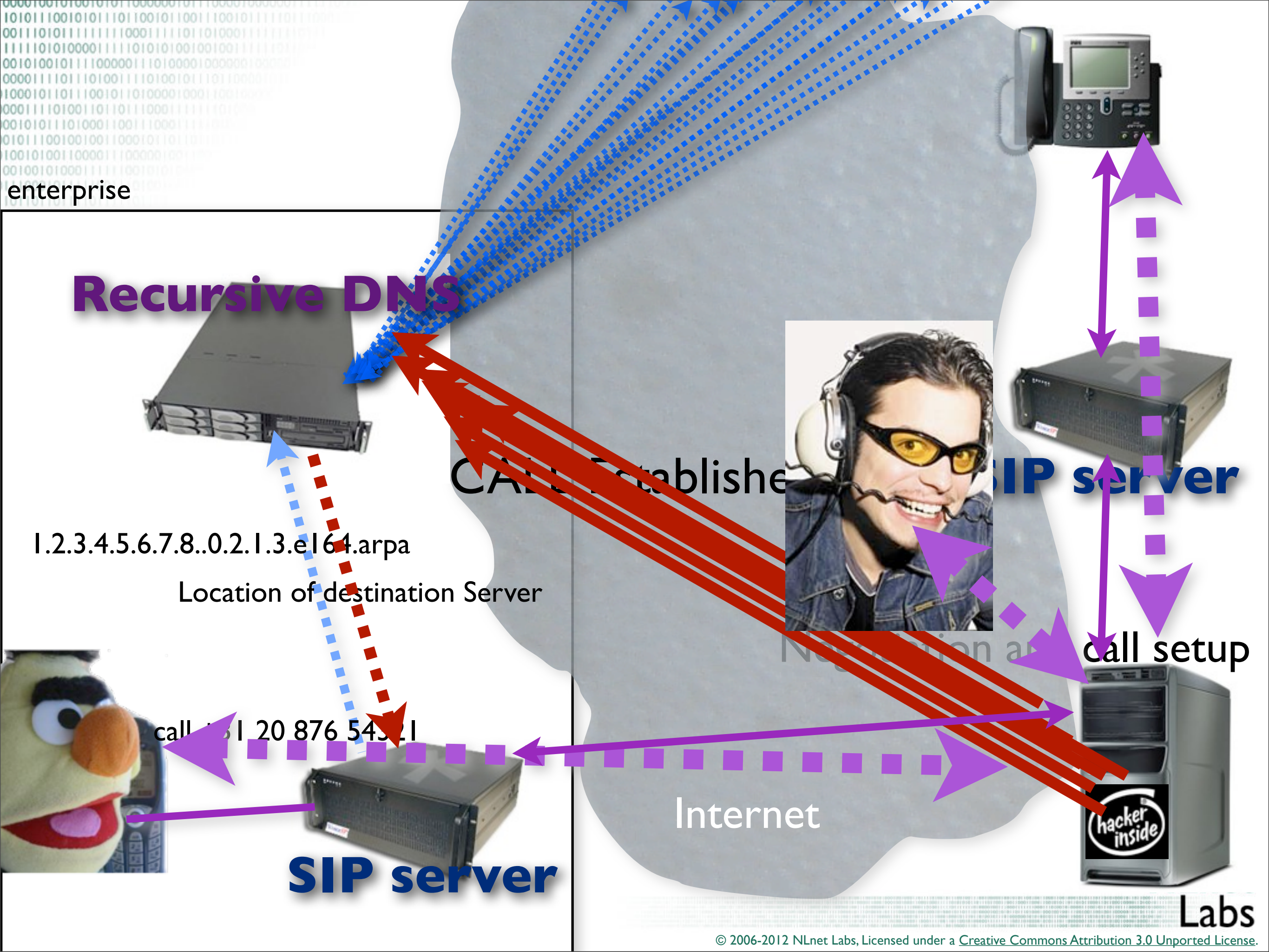
- When downloading Software upgrades
- Sharing their agenda
- Uploading tax forms
- Instant messaging with friends
- Connect to their security camera
- Figure out the latest news about that merger

# So DNS is IMPORTANT

- How would an attacker use the DNS for attacks?
- By fooling the receiver that a service lies elsewhere

Back to our VOIP example





# Cache Poisoning

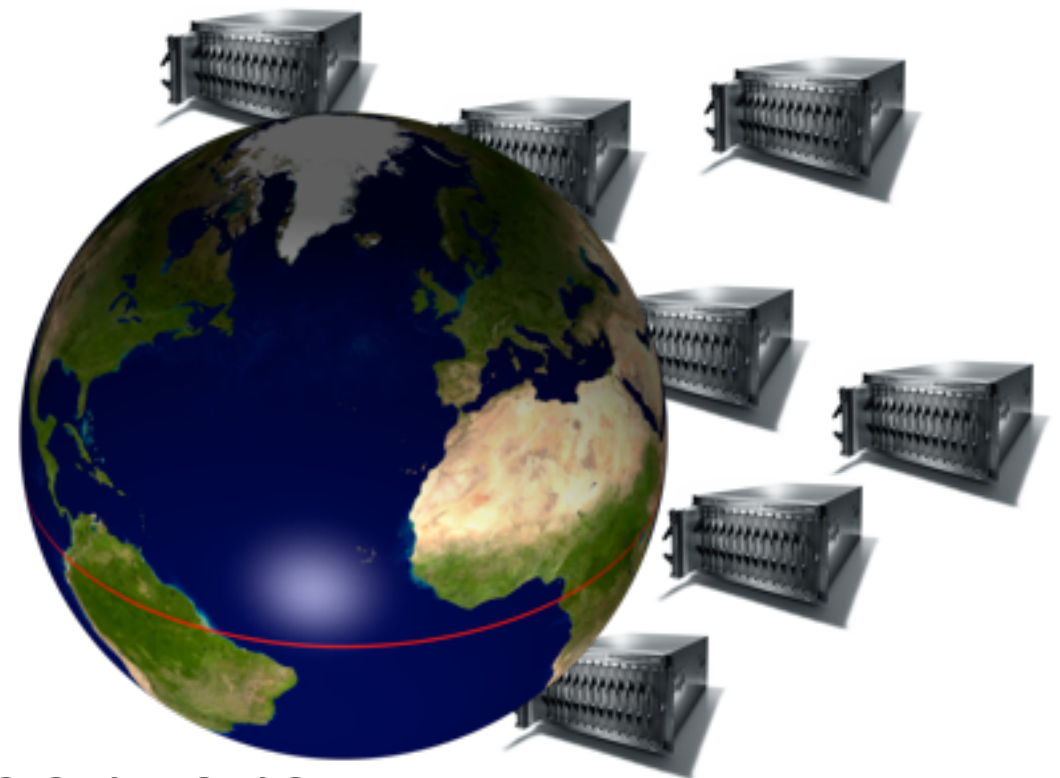
- The attack you just saw is called cache poisoning
- Inserting false data into the cache of recursive name servers
- This form of attack has been known for years
- One of the reasons to work on DNSSEC



# DNS Architecture and Protocol



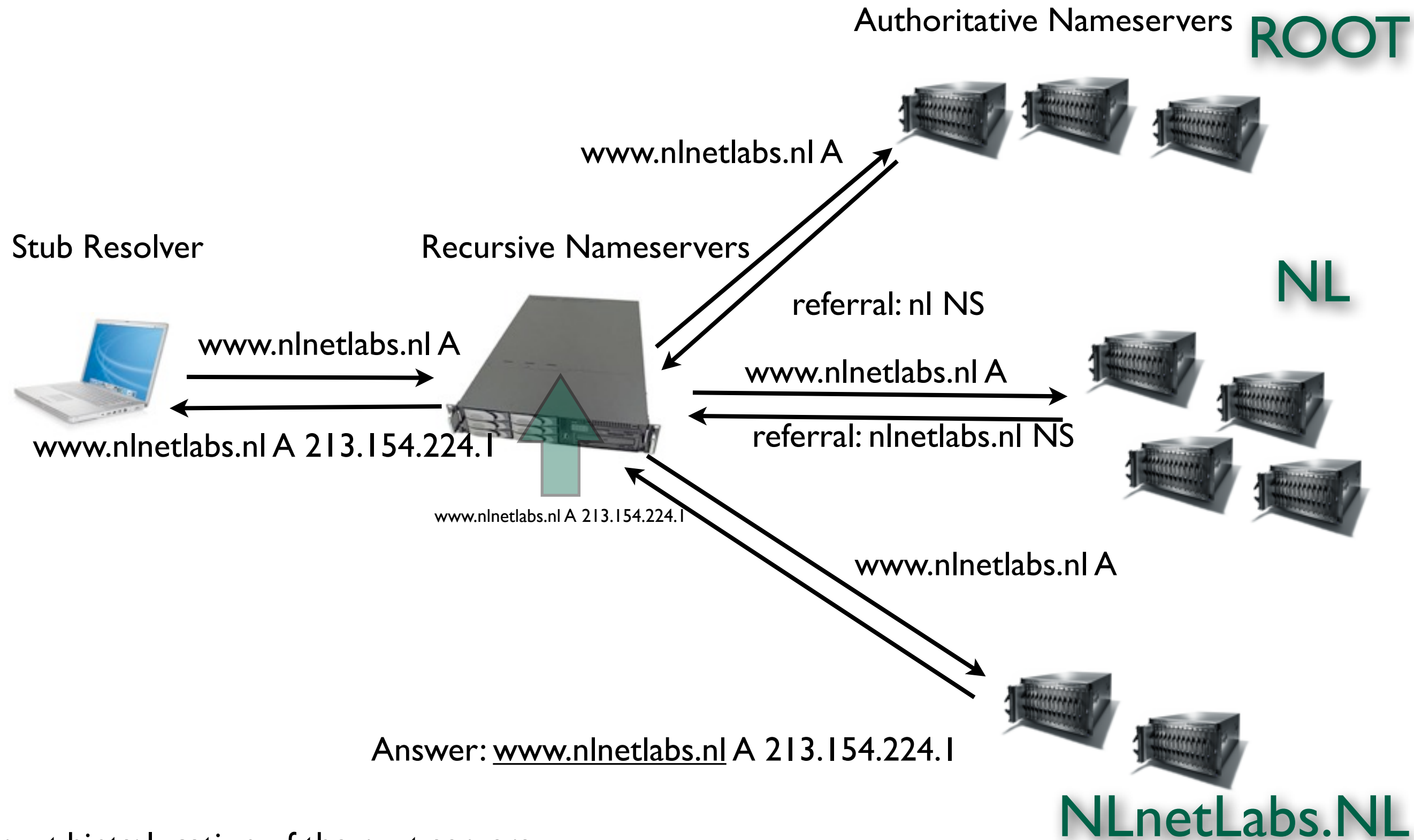
# Authoritative Nameservers



Stub Resolver

Recursive Nameservers





root.hints: location of the root servers

```
; <<>> DiG 9.7.0b2 <<>> @k.root-servers.net www.nlnetlabs.nl
; (2 servers found)
;; global options: +cmd
;; Got answer:
;; ->>HEADER<- opcode: QUERY, status: NOERROR, id: 41886
;; flags: qr rd; QUERY: 1, ANSWER: 0, AUTHORITY: 7, ADDITIONAL: 12
;; WARNING: recursion requested but not available
```

```
;; QUESTION SECTION:
;www.nlnetlabs.nl. IN A
```

```
;; AUTHORITY SECTION:
```

```
nl.      172800 IN NS nl1.dnsnode.net.
nl.      172800 IN NS ns1.nic.nl.
nl.      172800 IN NS ns2.nic.nl.
nl.      172800 IN NS ns3.nic.nl.
nl.      172800 IN NS ns4.nic.nl.
nl.      172800 IN NS ns-nl.nic.fr.
nl.      172800 IN NS sns-pb.isc.org.
```

```
;; ADDITIONAL SECTION:
```

```
nl1.dnsnode.net. 172800 IN A 194.146.106.42
ns1.nic.nl.      172800 IN A 193.176.144.2
ns2.nic.nl.      172800 IN A 213.154.241.28
ns3.nic.nl.      172800 IN A 194.171.17.2
ns4.nic.nl.      172800 IN A 62.4.86.232
ns-nl.nic.fr.    172800 IN A 192.93.0.4
sns-pb.isc.org.  172800 IN A 192.5.4.1
ns1.nic.nl.      172800 IN AAAA 2a00:d78::102:193:176:144:2
ns2.nic.nl.      172800 IN AAAA 2001:7b8:606::28
ns3.nic.nl.      172800 IN AAAA 2001:610:0:800d::2
ns-nl.nic.fr.    172800 IN AAAA 2001:660:3005:1::1:2
sns-pb.isc.org.  172800 IN AAAA 2001:500:2e::1
```

```
;; Query time: 4 msec
;; SERVER: 2001:7fd::1#53(2001:7fd::1)
;; WHEN: Tue Apr 6 14:12:44 2010
;; MSG SIZE rcvd: 447
```

Question

Referral

# Cache and TTL

;; ANSWER SECTION:

www.nlnetlabs.nl. 10200 IN A 213.154.224.1

- TTL is a parameter that indicates how long data is to remain in a cache
- TTL value is set by the zone owner
- TTL decreases while in the cache

# Back to Cache Poisoning



# Cache Poison

- Attack is based on ‘predicting’ properties
  - e.g. when asking a question to a female you expect a female voice to answer
- If you ask a question with a specific QID you expect that QID in the answer
- Cache poisoner will take a wild guess



```
; <<>> DiG 9.7.0b2 <<>> @k.root-servers.net www.nlnetlabs.nl
; (2 servers found)
;; global options: +cmd
;; Got answer:
;; ->>HEADER<- opcode: QUERY, status: NOERROR, id: 41886
;; flags: qr rd; QUERY: 1, ANSWER: 0, AUTHORITY: 7, ADDITIONAL: 12
;; WARNING: recursion requested but not available
```

```
;; QUESTION SECTION:
;www.nlnetlabs.nl. IN A
```

```
;; AUTHORITY SECTION:
nl.          172800 IN NS nl1.dnsnode.net.
nl.          172800 IN NS ns1.nic.nl.
nl.          172800 IN NS ns2.nic.nl.
nl.          172800 IN NS ns3.nic.nl.
nl.          172800 IN NS ns4.nic.nl.
nl.          172800 IN NS ns-nl.nic.fr.
nl.          172800 IN NS sns-pb.isc.org.
```

```
;; ADDITIONAL SECTION:
nl1.dnsnode.net. 172800 IN A 194.146.106.42
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ns3.nic.nl.      172800 IN A 194.171.17.2
ns4.nic.nl.      172800 IN A 62.4.86.232
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sns-pb.isc.org.  172800 IN AAAA 2001:500:2e::1
```

```
;; Query time: 4 msec
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;; WHEN: Tue Apr 6 14:12:44 2010
;; MSG SIZE rcvd: 447
```



# Varying properties in a packet

- The sender can vary the following properties for the attacker to match
- DNS:
  - Query ID (16 bits)
- Transport:
  - Fire the question from a random source port (16 bits)

# Isn't Query ID only sufficient?

Chance that  $n$  people have different birthdays

$$\bar{p}(n) = 1 \times \left(1 - \frac{1}{365}\right) \times \left(1 - \frac{2}{365}\right) \cdots \left(1 - \frac{n-1}{365}\right) = \frac{365 \times 364 \cdots (365 - n + 1)}{365^n} = \frac{365!}{365^n(365 - n)!}$$

Chance that  $n$  people have the same birthday

$$p(n) = 1 - \bar{p}(n).$$

n	P(n)
10	11.17%
20	41.1%
23	50.7%
30	70.6%
50	97%
57	99%
100	99.99997%

Bits	50%	5%	Aka
16	10 s	1 s	Unpatched server, random ID
26	2.8 h	17 m	Patched, using only 1024 ports
34	28 days	2.8 days	unbound with defaults
44	28444 days	2844.4 days	unbound with 0x20 and source addresses configured

# 50%-5%-0.5%-0.05%

# Besides: randomness is non-trivial

- For example: BIND9.4.1 and earlier used a pseudo random number generator that provided predictable sequences
- Current ID even: next ID one out of 10 possible numbers
- Only order 15 queries needed to predict rest of the stream
- Discovered by Amit Klein of trusteeer



# Using all ports, not easy

- Some architectures did not use a sufficiently large range of ports
- The patches issued as response to the so called Kaminsky attack, early 2008, all had to do with increasing the randomness in port use

# Still until 2007 folk seemed happy

- Attacker only got one try:
  - Query for `www.onlinebank.example`
  - Bombard with answers hoping for the the mala-fide answer to get in first
  - Wait for timeout of the TTL
  - Then try again

# Kaminsky's variant

- Classic cache poisoning gave you 'a few tries' to get in between the outgoing question and incoming answer
- Kaminsky came with a scheme where the culprit can keep trying
- Surprisingly simple, a wonder nobody thought of the variety before

# And how does it work

- Attacker queries:  
<randomcruft>.www.importantbank.example
- respond with fake delegation to:  
www.importantbank.example with glue
- There are other varieties too, but this is the one that has no real workaround

10101110010101110110010110011100101111011001100110000010111000010000001111  
001110101111111100011110110100001111110111  
111110101000011110101010010010011111011001  
001010010111000001101000010000001000001  
00001110111010011101001011101100001111  
1000101101110010110100001000110010001  
00011110100110110111000111111010101  
0010101110100011001110001111010101  
010111001001001100010110110110111  
100101001100001110000010011001  
001001010001111100101010101  
11100010111100111010011111  
1011011011110111101111  
00010110010100101001  
100011100100100101  
1110110111001111  
1100110000011111  
10111100001111  
011010101111  
010101011111  
111111  
111111  
111111

# problem?

# There is Recognition



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## Vulnerability Note VU#800113

### Multiple DNS implementations vulnerable to cache poisoning

#### Overview

Deficiencies in the DNS protocol and common DNS implementations facilitate DNS cache poisoning attacks.

#### I. Description

The Domain Name System (DNS) is responsible for translating host names to IP addresses (and vice versa) and is critical for the normal operation of internet-con

<http://www.kb.cert.org/vuls/id/800113>

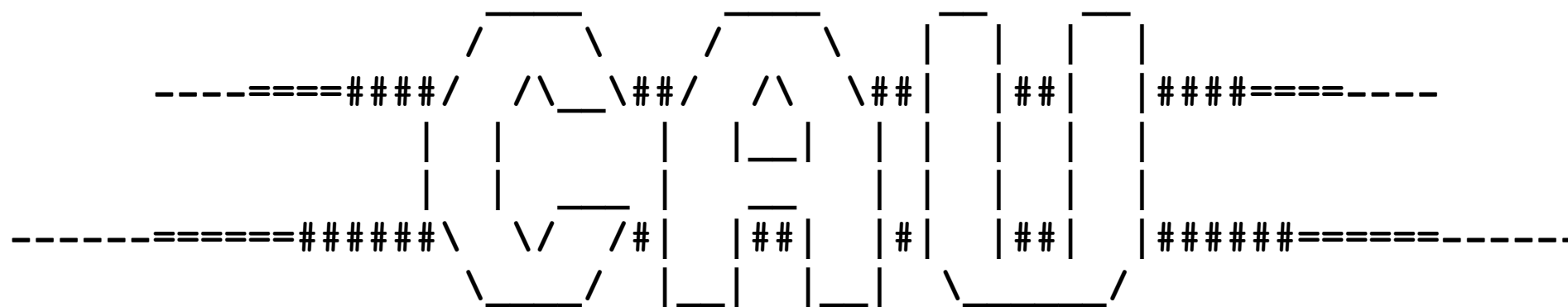
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# There is Exploit Code



Computer Academic Underground

<http://www.caughq.org>

Exploit Code

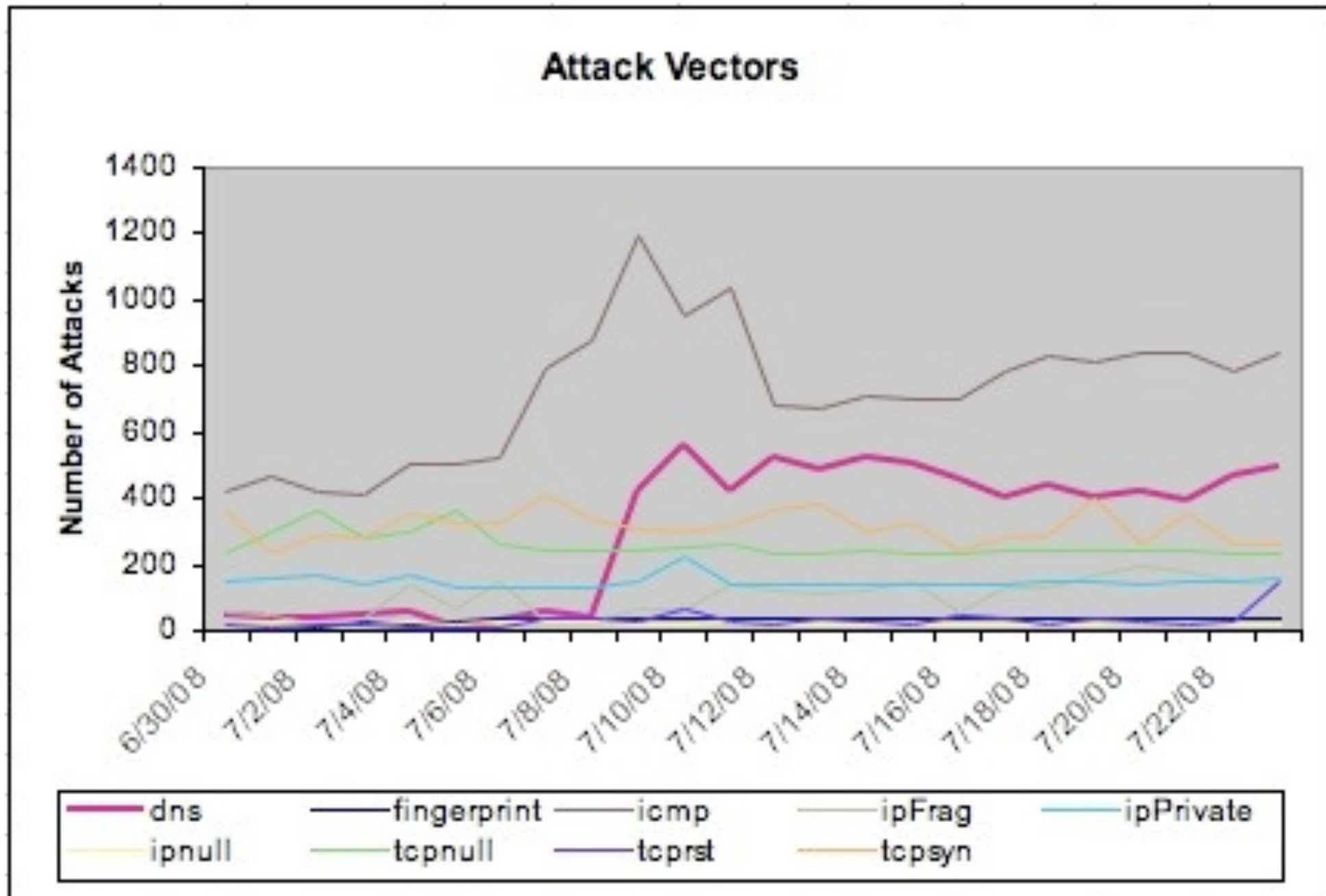
```
=====/  
Exploit ID:    CAU-EX-2008-0002  
Release Date:  2008.07.23  
Title:        bailiwicked_host.rb  
Description:   Kaminsky DNS Cache Poisoning Flaw Exploit  
Tested:       BIND 9.4.1-9.4.2  
Attributes:    Remote, Poison, Resolver, Metasploit  
Exploit URL:   http://www.caughq.org/exploits/CAU-EX-2008-0002.txt  
Author/Email:  I)ruid <druid (@) caughtq.org>  
              H D Moore <hdm (@) metasploit.com>  
=====/
```

# And more exploit code

```
/*  
* 2008+ Copyright (c) Evgeniy Polyakov <johnpol@2ka.mipt.ru>  
* All rights reserved.  
*  
* This program is free software; you can redistribute it and/or modify  
* it under the terms of the GNU General Public License as published by  
* the Free Software Foundation; either version 2 of the License, or  
* (at your option) any later version.  
*  
* This program is distributed in the hope that it will be useful,  
* but WITHOUT ANY WARRANTY; without even the implied warranty of  
* MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. See the  
* GNU General Public License for more details.  
*/
```


<http://tservice.net.ru/~s0mbre/archive/dns/>

# The networks are scanned



<http://asert.arbornetworks.com/2008/07/30-day-of-dns-attack-activity/>

# There have been successful attacks

**Today's Internet Threat Level: GREEN**  
Handler on Duty: Jim Clausung

GREEN

Diary Trends Reports About Presentations Top 10 Contact

Handler's Diary: Joomla user password reset vulnerability being actively exploited;Upcomi

## Diary

[previous](#) [next](#)

### DNS Cache Poisoning Issue Update

Published: 2008-07-30,  
Last Updated: 2008-07-30 21:20:49 UTC  
by David Goldsmith (Version: 1)

4 comment(s) [Digg](#) [submit](#)

Ok, we have a confirmed instance where the DNS cache poisoning vulnerability was used to compromise a DNS server belonging to AT&T. This PCWorld [article](#) covers the incident. The original article makes it sound as though the Metasploit site was 'owned' by this incident when really the issue was that the AT&T DNS server was compromised and was providing erroneous IP addresses to incoming queries. This updated PCWorld [article](#) clarifies the first one.

Additional details can be found in this Metasploit [blog post](#).

So we've moved from "the bad guys are out there" past "the invaders are at the gate" and on to "the bad guys are slipping inside". If your organization has not yet patched your DNS servers (see [here](#)), please do so now.

We may be raising our InfoSec status to yellow soon to help raise attention to the serious nature of this issue.

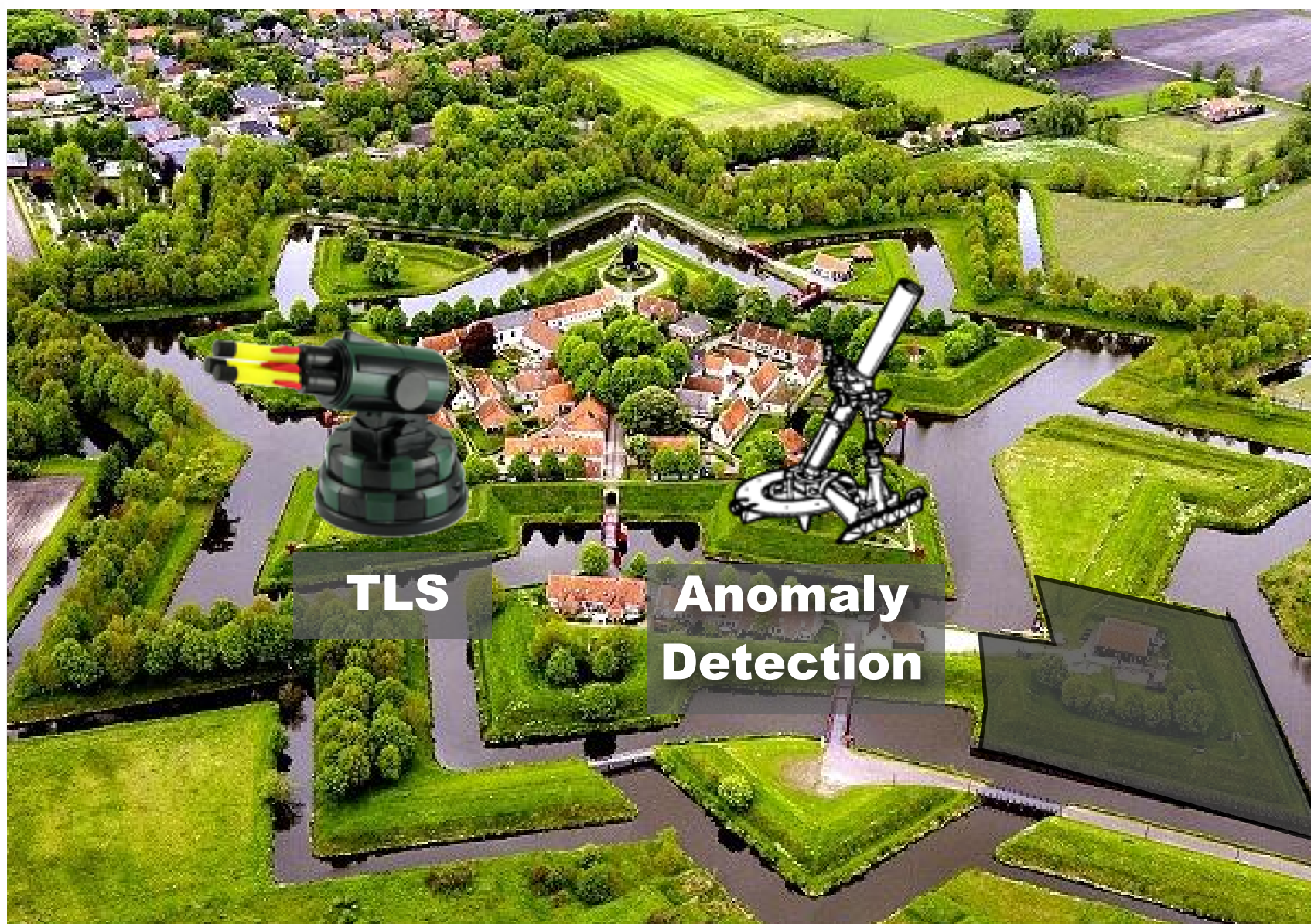
<http://isc.sans.org/diary.html?storyid=4801>



10101110010101110110010110011100101111011001111  
0011101011111110001111011010001111110111  
111110101000011110101010010010011111011011  
0010100101110000011101000010000001000001  
00001110111010011101001011101100001111  
1000101101110010110100001000110010001  
0001111010011011011100011111101011  
00101011101000110011100011110111  
0101110010010011000101101101111  
10010100110000111000001001100  
00100101000111110010101011  
1110001011110011101001111  
1011011011110111101111  
000101100101001011001  
100011100100100101  
1110110111001111  
1100110000011111  
10111100001111  
011010101111  
010110101111  
111111  
111111

# Yes, Problem





We lost DNS...

How about the other defenses ?



# SSL?

- Current practices are sloppy
- Users connect to their banks
- Get redirected to unrelated domains
- User interfaces only show padlocks

# For example

## Mastercard

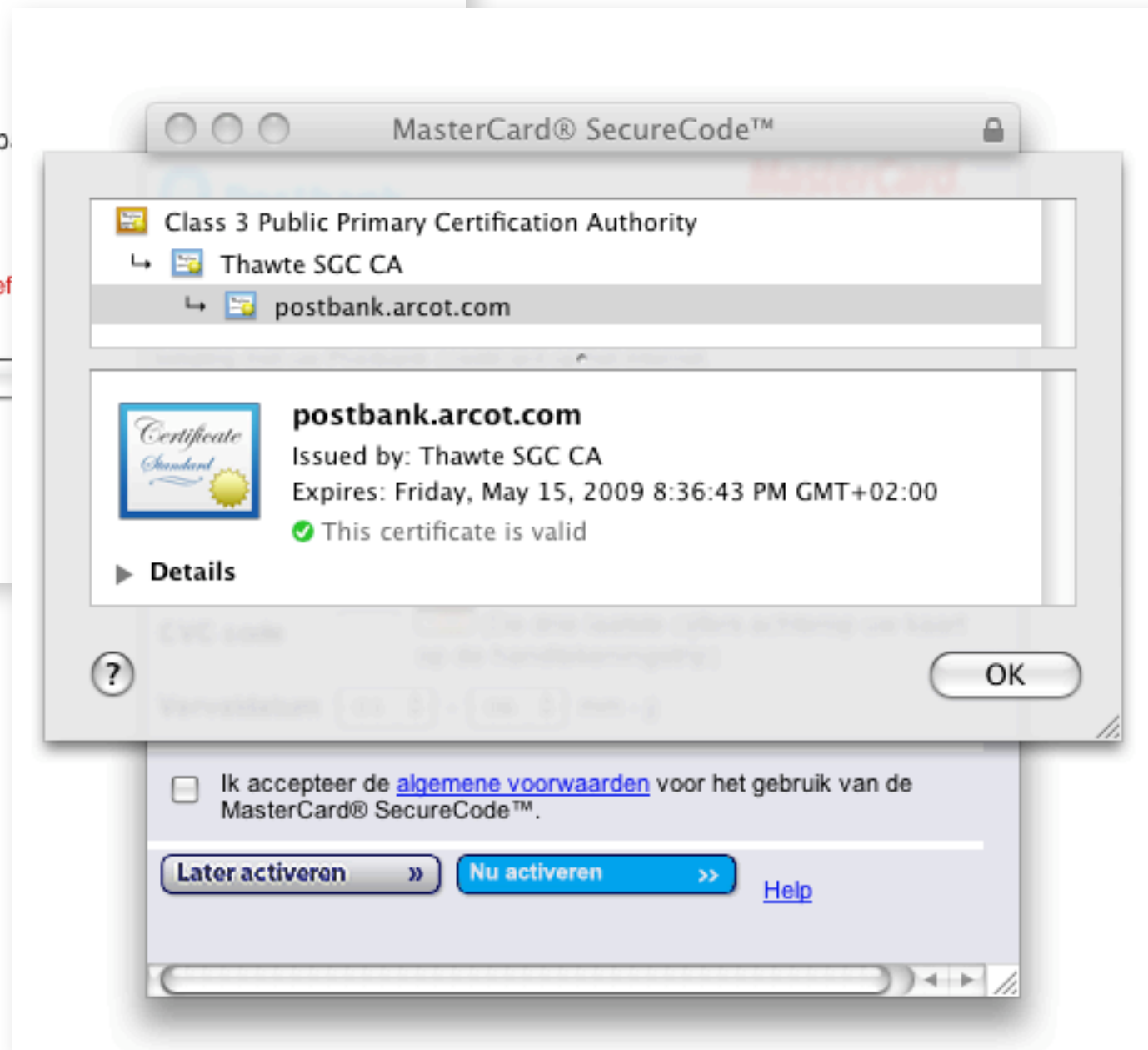


Press "Go to Your Bank" to authorise your credit card p

Amount **415.00 Euro**

Payment cluster ID **167102578**

Please deactivate your pop-up killer in your web browser, bef  
[more Info](#)



# Exploit

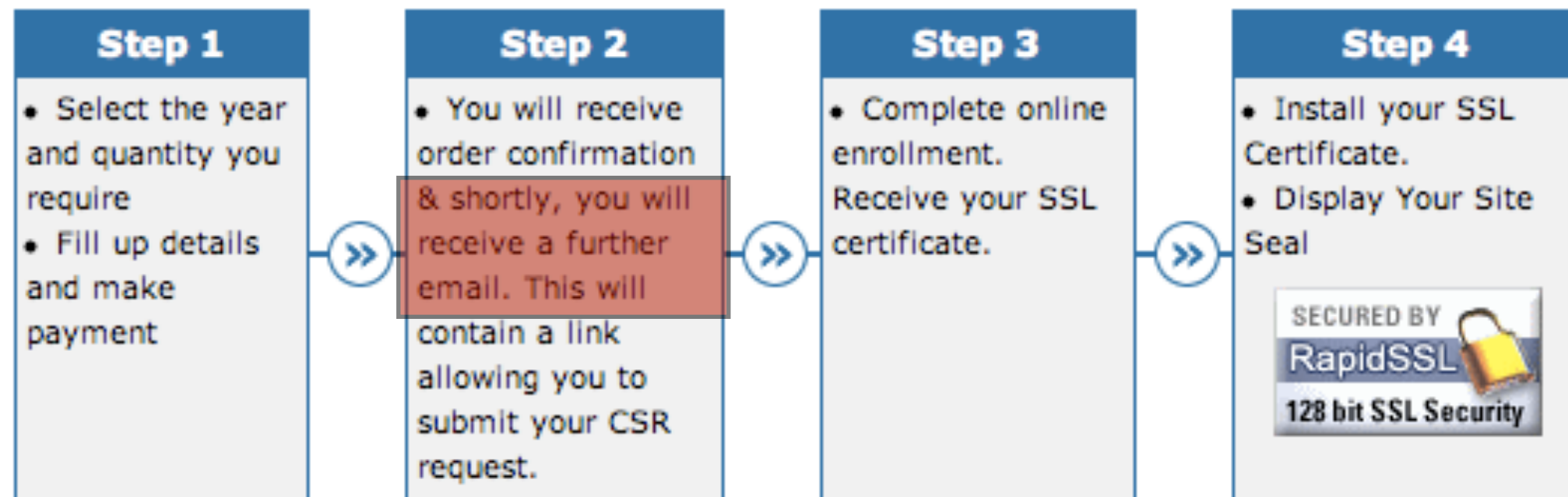
- Attacker poisons DNS for www.postbank.nl
- Fake www.postbank.nl redirects to postbank.webbanksecurity.com
- Obtaining the domain name and certificate is trivial for organized criminals
- Users are used to these sort of redirections and the domainname looks trustworthy

# Things get worse

- Fake [www.postbank.nl](http://www.postbank.nl) redirects to fake <https://www.postbank.nl>
- SSL protects against that?
- Not if the attacker has a signed certificate
  - How would an attacker do that?

## How SSL purchase works?

Ordering SSL from rapidsslonline.com online store is easy, fast and secure!  
You need to go through 4 simple steps to complete your SSL order



\*\*\* As part of GeoTrust's ongoing commitment to prevent fraud, some orders are randomly flagged for an additional security review. Please note that this order will not be fulfilled until GeoTrust completes this manual security review. Usually such orders are processed within 24 hours but sometimes may take longer than 24 hours. Please contact us via Email or Live Chat for Support in such cases.

<http://www.rapidsslonline.com/index.php>



# Don't rely on DNS for the Security review

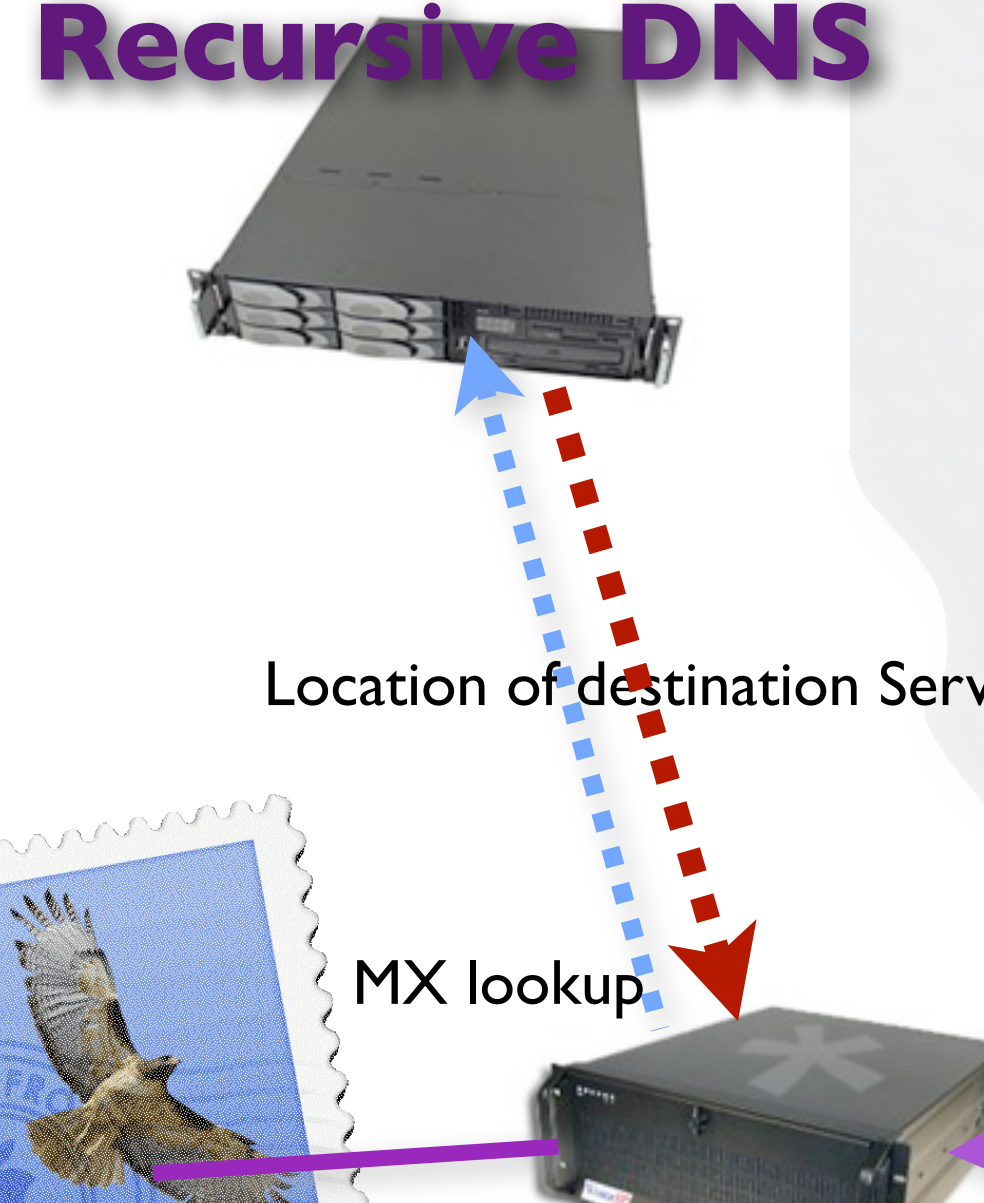
- Don't get the contact details out of the WHOIS, getting to WHOIS is DNS based
- Don't send confirmation e-mails to typical addresses in the domain
- Mail uses the DNS
- Don't try to see if domain already has a SSL certificate installed. That uses the DNS

# Lower hanging fruit: email

- Just attack e-mail
- Eavesdropping on e-mail
- Modifying text
- Inserting malicious content

enterprise

# Recursive DNS



The diagram illustrates the recursive DNS process. At the top, a rack server represents the Recursive DNS server. Below it, a smaller server represents the Mail server. A blue dashed arrow points from the Mail server up to the Recursive DNS server, labeled "MX lookup". A red dashed arrow points from the Recursive DNS server down to the Mail server, labeled "Location of destination Server". A purple arrow points from a postage stamp (featuring an eagle and an Apple logo) to the Mail server. The text "Mail server" is written in large blue letters at the bottom right.

Location of destination Server

MX lookup

Mail server

# DNSSEC

The diagram illustrates a mail server attack. At the top, a smartphone displays a 'Welcome to Mail' screen with logos for Y!MAIL, Gmail, .mac, AOL, and Other. A solid purple arrow points from the phone to a black server unit labeled 'Mail server'. A dashed purple arrow points from the 'Mail server' to a tower server unit labeled 'hacker inside'. A dashed purple arrow points from a black hat icon to the tower server. A long dashed purple arrow at the bottom is labeled 'Internet'.

# Labs

# Technique to notice these attacks

- SPF protocol for spam recognition
  - Based on... DNS
- TLS based connections and certification
  - In practice only used for encryption of the channel
  - Often misconfigured, or with fallback in place
  - And remember the problems wrt TLS



# EV vs DV

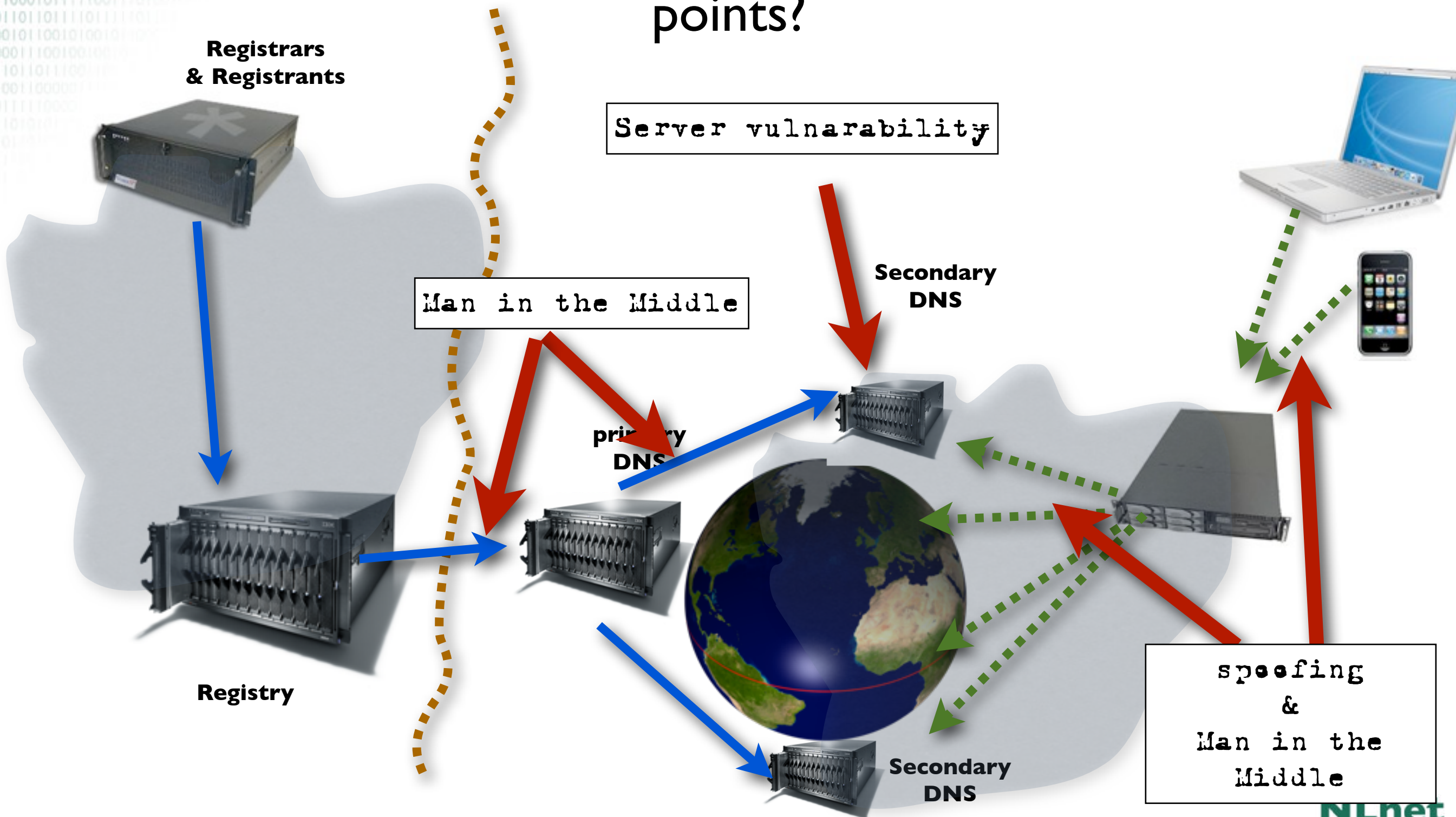
- Certificates come in two types: Domain Validation (DV) and Extended Validation (EV)
- Cert.Authorities hand out DVs purely on DNS based knowledge
- Difference: Green Glow in the browser (assuming UI is available)

[illegible]

# Is cache poisoning the only vulnerability?

# Data flow through the DNS

## Where are the vulnerable points?





# Protecting (Authoritative) Servers: Host Security

- Harden your OS
  - No unnecessary services/software
  - SSH with public keys only
  - Audit
- Harden your DNS secondary service provider
  - SLAs

# Protecting (Authoritative) Servers: Host Security II

- Run up-to-date software
  - OS stack
  - Nameserver software
- Software protection
  - chroot/jail environment
  - drop elevated permissions

# Protecting (Recursive) Servers

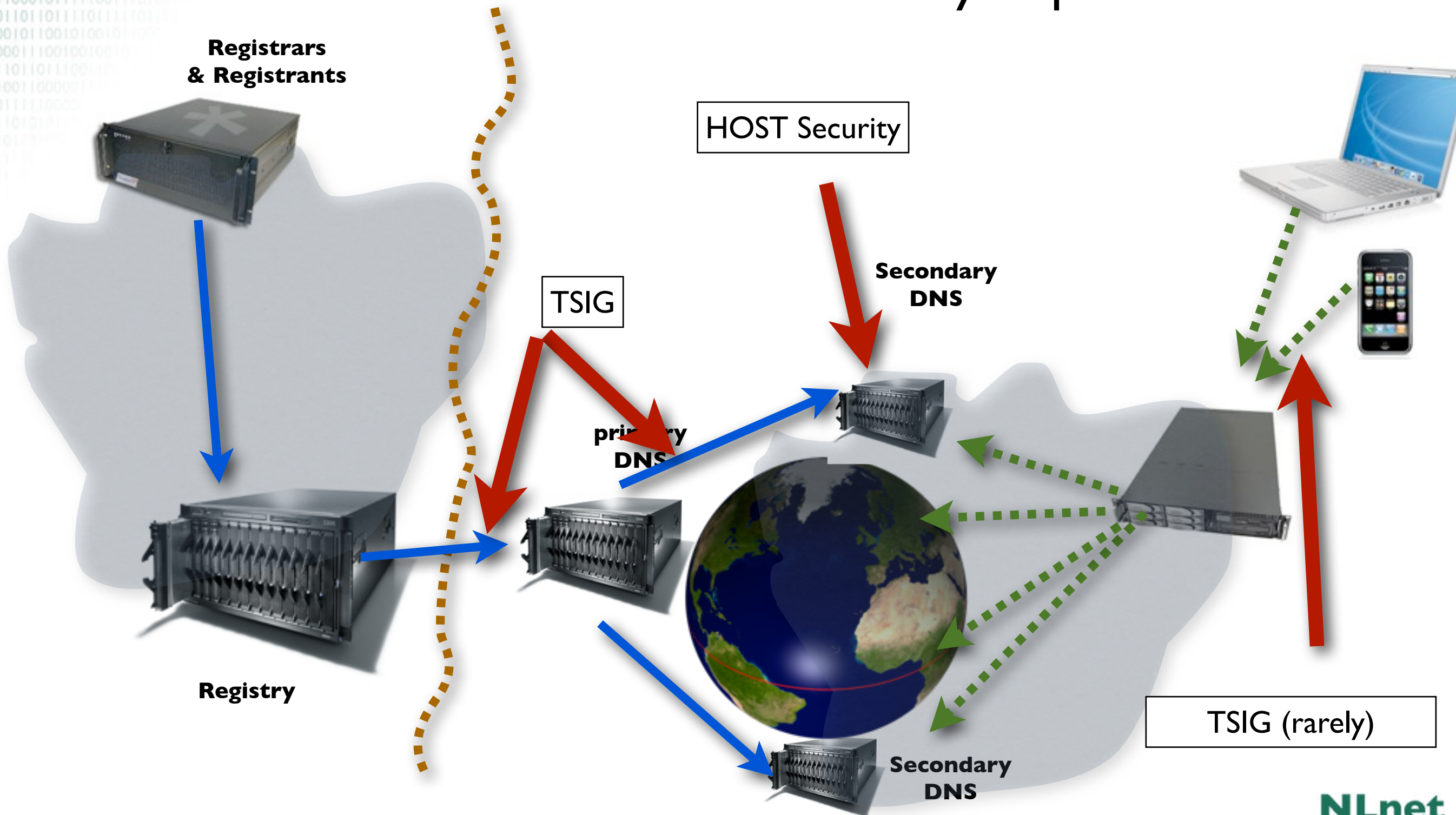
- Who do you accept questions from?
- Risks:
  - DOS by others (others may use your resources)
  - DOS to others (amplification attacks)
  - Have you implemented BCP 38?

# Securing Host-Host Communication



# Data flow through the DNS

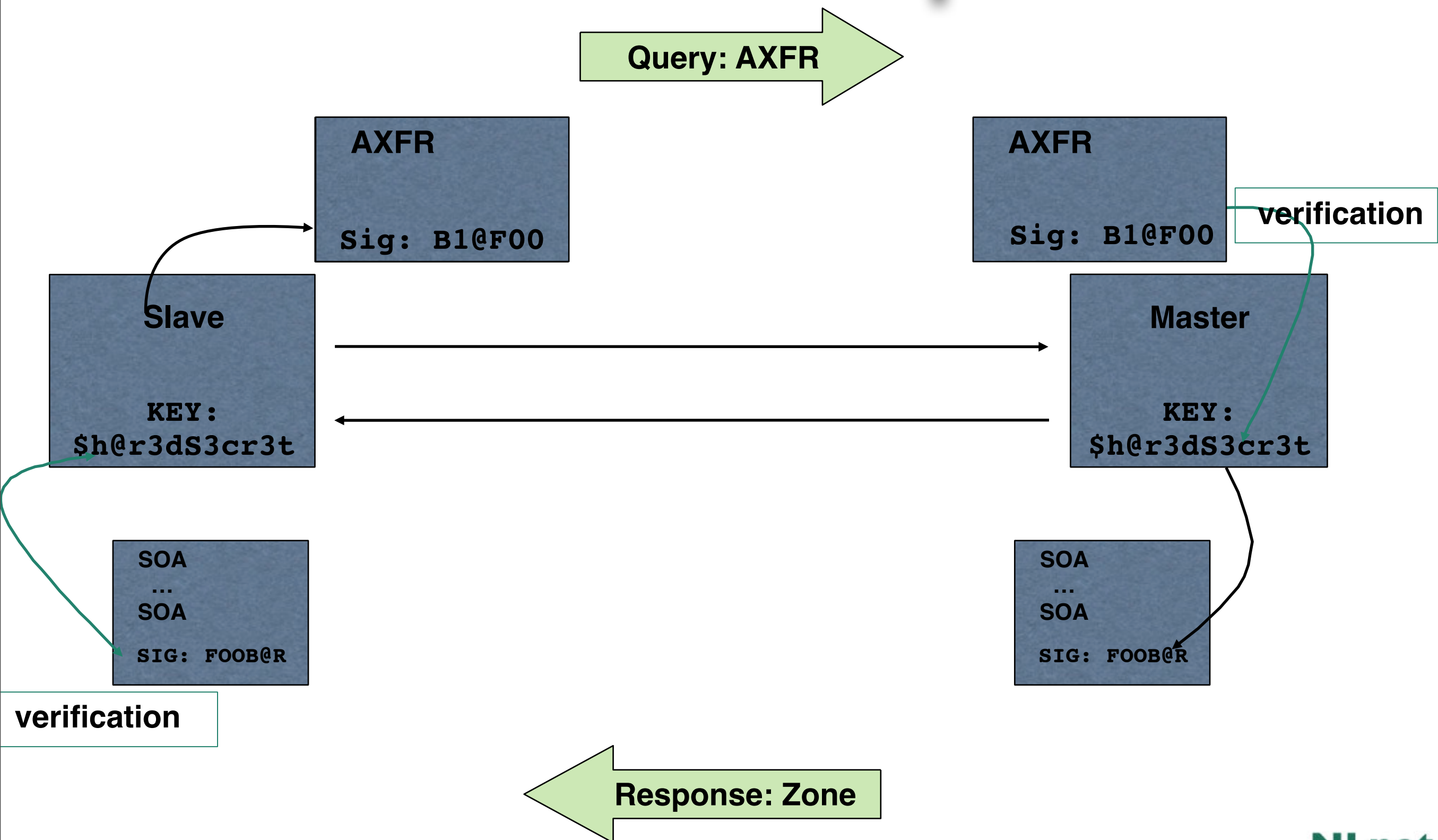
## What should you protect...



# Transaction Signature: TSIG

- TSIG (RFC 2845)
  - Authorising dynamic updates and zone transfers
  - Authentication of caching forwarders
  - Independent from other features of DNSSEC
- One-way hash function
  - DNS question or answer and timestamp
- Traffic signed with “shared secret” key
- Used in configuration, **NOT** in zone file

# TSIG Example



# TSIG for Zone Transfers

1. Generate secret
2. Communicate secret
3. Configure servers
4. Test



# Importance of the Time Stamp

- TSIG/SIG(0) signs a complete DNS request / response with time stamp
  - To prevent replay attacks
  - Currently hardcoded at five minutes
- Operational problems when comparing times
  - Make sure your local time zone is properly defined
  - `date -u` will give UTC time, easy to compare between the two systems
  - Use NTP synchronisation!

# Authenticating Servers Using SIG(0)

- Alternatively, it is possible to use SIG(0)
  - Not yet widely used
  - Works well in dynamic update environment
- Public key algorithm
  - Authentication against a public key published in the DNS
- SIG(0) specified in RFC 2931

# Cool Application

- Use TSIG-ed dynamic updates to configure your laptops name
- My laptop is know by the name of aagje.secret-wg.org
  - <http://ops.ietf.org/dns/dynupd/secure-ddns-howto.html>
  - Mac OS users: there is a bonjour based tool.
    - [www.dns-sd.org](http://www.dns-sd.org)

How about Unbound?

Unbound





# Security Choices in Unbound

- In general, a modern paranoid resolver
- DNSSEC support.
- RFC 2181 support completely
  - Fine grained. Keeps track of where RRSets came from and won't upgrade them into answers.
  - Does not allow RRSets to be overridden by lower level rrsets

# Filtering

- Scrubber:
- Only in-bailiwick data is accepted in the answer
  - The answer section must contain only answer
  - CNAME, DNAME checked that chain is correct
    - CNAME cut off and only the first CNAME kept
      - Lookup rest yourself do not trust other server
    - DNAME synthesize CNAME by unbound do not trust other server. Also cut off like above.
  - DNAME from cache only used if DNSSEC-secure.

# Filtering II

- No address records in authority, additional section unless relevant – i.e. mentioned in a NS record in the authority section.
- Irrelevant data is removed
  - When the message only had preliminary parsing and has not yet been copied to the working region of memory

# Entropy

- Randomness protects against spoof
  - Arc4random() (OpenBSD): crypto strong.  
May not be perfectly random, but predicting it is a cryptographical breakin.
  - Real entropy from OS as seed
- Query id – all 16 bits used.
- Port randomisation – uses all 16bits there, goes out of its way to make sure every query gets a fresh port number



# Entropy II

- Destination address, and ipv4/ipv6. RTT band of 400msec (=everything).
- Its not the timewindow but the randomness
- Query aggregation – same queries are not sent out – unless by different threads
- Qname strict match checked in reply
- 0x20 option
- Harden-referral-path (my draft) option
- Can use multiple source interfaces!
  - 4 outgoing IP address add +2 bits



# Other measures

- Not for the wire itself
  - Heap function pointer protection (whitelisted)
  - Chroot() by default
  - User privileges are dropped (lots of code!)
  - ACL for recursion
  - No detection of attacks – assume always under attack
  - version.bind hostname.bind can be blocked or configured what to return (version hiding)
  - Disprefer recursion lame servers – they have a cache that can be poisoned

